Single-shot number-resolved detection of microwave photons with error mitigation

Jacob C. Curtis\textsuperscript{1,2,\dagger}, Connor T. Hann\textsuperscript{1,2}, Salvatore S. Elder\textsuperscript{1,2}, Christopher S. Wang\textsuperscript{1,2}, Luigi Frunzio\textsuperscript{1,2}, Liang Jiang\textsuperscript{3} and Robert J. Schoelkopf\textsuperscript{1,2,†}

\textsuperscript{1}Departments of Physics and Applied Physics, Yale University, New Haven, Connecticut 06511, USA
\textsuperscript{2}Yale Quantum Institute, Yale University, New Haven, Connecticut 06511, USA
\textsuperscript{3}Pritzker School of Molecular Engineering, University of Chicago, Chicago, Illinois 60637, USA

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Single-photon detectors are ubiquitous and integral components of photonic quantum cryptography, communication, and computation. Many applications, however, require not only detecting the presence of any photons, but distinguishing the number present with a single shot. Here, we implement a single-shot, high-fidelity photon number-resolving detector of up to 15 microwave photons in a cavity-qubit circuit QED platform. This detector functions by measuring a series of generalized parity operators, which make up the bits in the binary decomposition of the photon number. Our protocol consists of successive, independent measurements of each bit by entangling the ancilla with the cavity, then reading out and resetting the ancilla. Photon loss and ancilla readout errors can flip one or more bits, causing nontrivial errors in the outcome, but these errors have a traceable form, which can be captured in a simple hidden Markov model. Relying on the independence of each bit measurement, we mitigate biases in ensembles of measurements, showing good agreement with the predictions of the model. The mitigation improves the average total variation distance error of Fock states from 13.5%–1.1%. We also show that the mitigation is efficiently scalable to an M-mode system provided the errors are independent and sufficiently small. Our work motivates the development of new algorithms that utilize single-shot, high-fidelity PNR detectors.

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I. INTRODUCTION

Many quantum information processing protocols that utilize photonic platforms require devices that detect single photons. These single-photon detectors (SPDs) perform essential tasks such as measuring the output of a photonic register or heralding in probabilistic quantum gates. Uses of SPDs include quantum key distribution [1–3], linear optical register or heralding in probabilistic quantum gates. Uses of essential tasks such as measuring the output of a photonic photons. These single-photon detectors (SPDs) perform es-

SPDs in multiplexing or arraying schemes [26–31] and transition edge sensors [32–34]. These recent advances have vastly improved performance, but the limited fidelity and resolvable number of optical PNR detectors hamper multiphoton experiments [24].

The detection of single microwave photons is less established and more challenging, due to their lower energy and higher thermal background levels. The dispersive interaction between a photon and an atom or qubit enables a wide range of measurement capabilities not possible in the optical domain. In cavity QED, the dispersive interaction with Rydberg atoms enabled the observation of single photon jumps [35]. In circuit QED (cQED) systems, Josephson junctions coupled to microwave cavities generate the dispersive interaction essential for measuring single microwave photons [36,37]. Additionally, the dispersive interaction has been used to create [38,39] and manipulate [40,41] these photons in a wide variety of ways. cQED analogs of optical photodetectors such as the Josephson photomultiplier [42–44] have also been developed.

Introducing high-performance PNR detectors to the microwave regime could greatly enhance the prospects for performing boson sampling and other photonic quantum information processing techniques with cQED systems. Early quantum nondemolition (QND), single-shot, number-resolving measurements of microwave photons used flying Rydberg atoms to iteratively update the experimentalists’ knowledge of the photon number distribution [45]. Previous efforts in cQED to develop PNR measurements have fallen short of combining single-shot and QND capabilities. Spectral
FIG. 1. Projective, binary-valued measurements. (a) Measuring a qubit in its energy eigenbasis (the basis of $\sigma_z$ eigenstates) projects the qubit into its ground state $|g\rangle$ with probability $|\alpha|^2$ or its excited state $|e\rangle$ with probability $|\beta|^2$. The measurement operators $\hat{M}_g$ and $\hat{M}_e$ commute with the Hamiltonian at all times. As a result, the shared eigenstates are not perturbed by additional measurements or time evolution, making the measurement QND. (b) Here, Fock states are represented with the first two bits (parity and superparity) of their binary decomposition $|n\rangle = |b_1(n)\rangle|b_0(n)\rangle$ for $N < 4$. An initial measurement of parity (black arrows labeled by measurement outcome and corresponding measurement operator) projects $|\psi\rangle$ into the even or odd subspace, and renormalizes the remaining state probabilities. A subsequent measurement of superparity (red arrows), further projects the state fully into a single Fock state $|n\rangle$. The detected photon number is computed using a record of each measurement.

density analysis [39] and state interrogation methods [36] are QND, but require many shots of the experiment to build up time or spectroscopic traces. Parity measurements have been used to count photon jumps [46], but, while single shot, the full PNR measurement is not QND. More recently, frequency multiplexing [47] and multiple feed forward measurements [48] have been used to make QND, single-shot PNR measurements, but suffer from limited fidelity.

In this paper, we describe a measurement protocol that implements a high-fidelity, single-shot PNR detector of microwave photons. The single-shot nature of this measurement is essential to sample from the exponentially large Hilbert space inherent to experiments with multiple bosonic modes. We previously introduced this protocol to efficiently sample from the probability distribution at the output of a multiphoton bosonic quantum simulator, reducing the required number of measurements by a factor of 256, the size of the Hilbert space [49]. Errors in the storage and ancilla modes caused nontrivial bias in the output, but these errors are well understood and can be modeled with a simple hidden Markov model. We calibrate and use this model to improve the fidelity of the measurement by an order of magnitude using error mitigation methods, reducing the measurement infidelity to a few percent.

This paper is structured as follows. First, in Sec. II, we discuss QND measurements of multiple observables with a single prepared state. We verify measurements of parity and its generalizations satisfy the conditions required to be QND. Next, we show how to construct a PNR detector of microwave photons from successive QND measurements of parity and its generalizations. This is accomplished by representing the photon number in its binary decomposition. In Sec. IV, we introduce the error syndromes of our system along with a hidden Markov model (HMM) that parametrizes the errors. After calibrating the error rates, we use the model to mitigate the errors on an ensemble of states using deconvolution methods. Finally, we show that our error deconvolution protocol is scalable in the number of modes and maintains computational efficiency.

II. BACKGROUND

Optical photon detectors such as photomultiplier tubes, avalanche photodiodes, and superconducting nanowire SPDs function by converting incoming photons into electrical charges, which are then amplified to produce a detectable signal [17]. This process is single shot, but is inherently destructive as the incoming photons are consumed by the detector. Standard techniques in cQED for measuring single photons are QND [37,46], allowing for additional processing of the state. Our protocol for PNR detection utilizes this ability to perform multiple measurements that collectively form a
number-resolving measurement. In this section, we generalize the notions of QND and single shot to a multilevel system with multiple measurements.

A. Multilevel QND measurements

To realize a single-shot PNR detector, we need to extend the concept of QND measurements to the case of a multilevel system and multiple measurements. Not all observables can be measured in a QND way; they must satisfy two conditions [50]. The first QND condition is that the observable $\hat{A}$ commutes with the system Hamiltonian $\hat{H}$ to ensure that a repeated measurement of it yields the same result at any subsequent time (in the absence of external perturbations, such as photon loss). The second condition is that $\hat{A}$ commutes with the measurement Hamiltonian $\hat{H}_M$, which contains additional terms that couple to a meter. This condition protects the state from additional evolution induced by the measurement process. QND measurements implement measurement operators $\hat{M}_\lambda$ for each outcome $\lambda$, which project the system into the eigenstate labeled by its eigenvalue $\lambda$.

Let us first consider the usual case of QND measurements of a qubit or two-level system initialized in an arbitrary state $|\psi\rangle = \alpha |g\rangle + \beta |e\rangle$. A measurement of the system in this basis, with measurement operators $\hat{M}_g = |g\rangle\langle g|$ and $\hat{M}_e = |e\rangle\langle e|$, projects $|\psi\rangle$ into $|g\rangle$ with probability $|\alpha|^2$ and $|e\rangle$ with probability $|\beta|^2$. The outcome of this measurement tells us exactly the state of the qubit following the measurement. If the measurement of the observable is QND, then repeated measurements at any subsequent time will yield the same result, as seen in Fig. 1(a).

How can we extend this concept to perform a QND measurement on a multilevel system? For instance, consider a four-level subspace of a bosonic mode $s$ with operators $\hat{\delta}$, $\hat{\delta}^\dagger$. Two bits of classical information must be obtained to determine the state of this system. This can be accomplished with two QND measurements, whose composition projects the system into a single level. The first QND measurement halves the Hilbert space, and projects the state into one of its two subspaces, extracting one bit of classical information. A second measurement further halves this subspace, extracting the second bit and projecting the state into a single level of $s$. This is only possible if the two operators commute, or else the second measurement would project into a superposition of states from the two subspaces of the first measurement, spoiling the first bit of information.

For example, the binary decomposition of the number of excitations in $s$ contains two bits identifying the state of the four-level system discussed above. Parity $\hat{P}^0$ divides the space into even and odd subspaces with measurement operators $\hat{P}^0_{b=0} = |0\rangle\langle 0| + |2\rangle\langle 2|$ and $\hat{P}^0_{b=1} = |1\rangle\langle 1| + |3\rangle\langle 3|$. This measurement projects the system into the state corresponding to measurement outcome $b_0$, which is the least significant bit in the binary decomposition of the excitation number of $s$. Measurements of superparity $\hat{P}^1$, which commutes with $\hat{P}^0$, measure the second-least-significant bit $b_1$ with measurement operators $\hat{P}^1_{b=0} = |0\rangle\langle 0| + |1\rangle\langle 1|$ and $\hat{P}^1_{b=1} = |2\rangle\langle 2| + |3\rangle\langle 3|$. Measured in succession, these two observables project the system into a single Fock state labeled by its binary decomposition $|b_1b_0\rangle$, as shown in Fig. 1(b). These generalized parity measurements form the basis of our approach to realize a QND, PNR detector of microwave photons in a cQED system.

B. Multilevel single-shot measurements

As we discuss in the previous section, measuring a two-level or multilevel system requires extracting one or more bits of information. For these measurements to be considered single-shot, they must accurately measure a significant majority of their bits in a single shot. When measuring a qubit, this means that one gains close to one bit of information about the state per measurement, per initial state subjected to that measurement. With a correspondingly high measurement efficiency and a low dark count rate, a qubit measurement then extracts enough of the available information that we can measure it with precision close to the shot noise limit imposed by the ensemble size.

We call a multilevel measurement single shot if a significant fraction of the bits are faithfully extracted per shot. For example, choosing $3/4$ as this fraction corresponds to extracting more than three bits of information per shot for a PNR detector with $N = 16$ levels. To determine $P(n)$ of a multilevel system, we must sample many times from an ensemble of states. If the PNR measurement that acquires each sample meets this threshold, the error in the sampled $P(n)$ is limited by the shot noise of the number of samples. This capability is particularly crucial to extract joint probability distributions from $M N$-level systems without individually checking each of the $N^M$ states. We discuss the single-shot character of our detector in Sec. IV B.

III. BITWISE MEASUREMENT

A. Measuring binary-valued cavity observables

Our PNR detector uses a qubit ancilla to measure binary-valued observables of a bosonic storage mode via the dispersive interaction. These observables $\hat{O}$ divide the Hilbert space into exactly two subspaces with distinct eigenvalues $\lambda_1$, $\lambda_2$. We can imagine a gate that entangles the two ancilla states with the eigenspaces $E_{\lambda_1}, E_{\lambda_2}$. A readout of the ancilla would then project the storage mode into the corresponding eigenspace. This two-step measurement is QND only if the QND conditions are satisfied for all steps in the measurement. For a measurement of $\hat{O}$ to satisfy the first QND condition, it must be diagonal in the energy eigenbasis.

We next verify that the second QND condition is satisfied. Our three-mode cQED system consisting of a readout mode, a storage mode, and an ancilla has operators $\hat{r}$, $\hat{p}^\dagger$, $\hat{s}$, $\hat{s}^\dagger$, and $\hat{\delta}$, respectively. The composite system Hamiltonian is

$$\hat{H}/\hbar = \frac{\omega_0}{2}\hat{\delta} + \omega\hat{s}\hat{s}^\dagger + \omega\hat{p}\hat{p}^\dagger - \frac{X_{\lambda_1}}{2}\hat{s}\hat{s}^\dagger - \frac{X_{\lambda_2}}{2}\hat{p}\hat{p}^\dagger,$$

(1)

The device is the same as used in Refs. [51,52] with similar parameters, unless otherwise noted. We first apply an ancilla-storage entangling drive, which produces the mapping Hamiltonian $\hat{H}_{\text{map}}/\hbar = \hat{H}/\hbar + \Omega(t)\hat{\delta}$. For observables $\hat{O}$, frequency-selective drives $\Omega(t)$ can implement the requisite entangling operation by exciting the ancilla conditioned on a Fock state. The second step, ancilla readout, also...
utilizes a dispersive interaction, realized with the Hamiltonian $$\hat{H}\text{readout} = \hat{H} + \epsilon (\hat{r} + e^{i\theta} \hat{r})$$. This readout completes the measurement of $$\hat{O}$$. Both the mapping and readout Hamiltonians commute with the diagonal $$\hat{O}$$, and satisfy the second QND condition. These drives can induce dephasing [52], but do not induce additional decay in the storage mode [46]. Measurements of $$\hat{O}$$ are thus QND, enabling the measurement of multiple generalized parity operators that form the basis of our number-resolving measurement.

### B. Generalized parity measurements

To implement the number-resolving measurement, we synthesize gates that map the binary-valued generalized parity operators onto the ancilla qubit. Each gate enables the measurement of one bit, so to resolve the first $$N$$ states, we require $$B = \log_2 N$$ gates. The generalized parity operators

$$\hat{P}^k_{ij} = \begin{cases} 0 & \text{for } i \neq j \\ 1 - 2|\frac{i}{2}|(\text{mod 2}) & \text{for } i = j \end{cases}$$

(2)

with eigenvalues $$\lambda = \pm 1$$ halve the Hilbert space into two eigenspaces $$E_{\pm} = |n| b_k(n) = 0$$ and $$E_{\mp} = |n| b_k(n) = 1$$, where $$b_k(n)$$ is the $$k$$th bit in the binary decomposition of $$n$$. The $$\hat{P}^k$$ are diagonal in the energy eigenbasis and thus satisfy both conditions for QND measurement. To perform this measurement, we use numerical optimal control techniques (see Appendix D for more details) to synthesize a CNOT-like unitary operation

$$\hat{C}^k \hat{P}^k = \sum_{b_k(n)=0} e^{i\phi_k} |n\rangle \langle n| \otimes |g\rangle \langle g| + \sum_{b_k(n)=1} e^{i\phi_k} |n\rangle \langle n| \otimes |e\rangle \langle e|$$

$$+ \sum_{b_k(n)=1} e^{i\phi_k} |n\rangle \otimes |g\rangle \langle e|.$$  

(3)

The gate imparts phases $$\phi_k, \phi_k$$, but these do not affect the final measurement outcome because the composite measurement projects into a single Fock state, rendering the induced phase irrelevant. We can set these phases ourselves, but here allow the optimizer flexibility to choose the phases. Applying this gate to an ancilla prepared in $$|g\rangle$$ coupled to a storage mode with state $$|\phi\rangle = \sum_n c_n |n\rangle$$ entangles the odd and even subspaces with the state of the ancilla,

$$\hat{C}^k \hat{P}^k (|\phi\rangle \otimes |g\rangle)$$

$$= \sum_{b_k(n)=0} c_n e^{i\phi_k} |n\rangle \otimes |g\rangle + \sum_{b_k(n)=1} c_n e^{i\phi_k} |n\rangle \otimes |e\rangle.$$  

(4)

QND measurements of the ancilla then project the storage state into either $$E_{\pm}$$ or $$E_{\mp}$$ depending on the outcome, realizing the measurement operators

$$\hat{B}^{(k)}_{(0)} = \sum_{b_k(n)=0|0(1)} e^{i\phi_k} |n\rangle \langle n|.$$  

(5)

The act of measuring $$B$$ parity operators $$\hat{P}^{B-1}, \ldots, \hat{P}^1, \hat{P}^0$$ with outcomes $$b_{B-1}, \ldots, b_1, b_0$$ in the subspace including only the first $$N$$ states is equivalent to applying the measurement operator

$$\hat{M}_n = \hat{B}_{b_{B-1}}^{(B-1)} \ldots \hat{B}_{b_1}^{(1)} \hat{B}_{b_0}^{(0)} = |n\rangle \langle n|$$  

(6)

and renormalizing. This operator projects the system into the Fock state $$|n\rangle = |b_{B-1}(n) \ldots b_1(n) b_0(n)\rangle$$ completing the bitwise measurement. This approach requires the minimal number of binary measurements of any scheme assuming no prior knowledge of the state. This protocol implements a detector resolving up to $$2^B$$ photons, where the number of generalized parity (bit) measurements we make is fully programmable.

### C. Experimental implementation

To resolve the storage mode photon number in a single shot, we sequentially measure the generalized parity operators, whose outcomes form the binary decomposition of the photon number. For each bit, we apply $$\hat{C}^k \hat{P}^k$$, read out the ancilla, and dynamically reset it using the same method as in Ref. [51]. More details of the reset protocol, including statistics on the number of attempts required, are in Appendix A. This process is concatenated to measure all four bits, allowing us to resolve up to 15 photons, as shown in Fig. 2(a).

Another approach performs a series of paritylike measurements, feeding forward the result of each bit measurement to determine the correct qubit rotation angle for the next measurement [48]. However, any errors that occur during a single measurement are fed forward and corrupt any subsequent measurements, resulting in correlated errors. Our method, which independently measures each bit, is more suitable for error mitigation with postprocessing techniques.

Each component of our measurement chain is susceptible to errors. The storage mode suffers from stochastic photon loss at a rate $$\kappa$$. The ancilla is vulnerable to dephasing and excitation decay during the mapping pulse $$\hat{C}^k \hat{P}^k$$ as well as decay during the readout procedure. The probability of reading out the ancilla in the ground (excited) state when the readout should have yielded $$e(g)$$ is $$\epsilon_{e(g)}$$. The long lifetimes of the storage mode $$T^g_i \approx 1$$ ms and ancilla $$T^g_i \approx 25 \mu s$$ relative to the 2.9 $$\mu s$$ duration of each measurement ensure that these errors are small. The dynamic reset protocol rarely leaves the ancilla in the excited state, and is not a significant source of error.

Counterintuitively, a single error in the entire sequence can produce a result that differs from the correct value by up to eight photons. This happens because a single error can flip multiple bits in the photon number’s binary decomposition. For example, the loss of a single photon from $$|8\rangle = |1000\rangle$$ after the measurement of $$b_0 = 0, b_1 = 0, b_2 = 0$$ results in $$b_3(7) = 0$$. Thus, the effect of such an error is to mistakenly read out $$|8\rangle = |1000\rangle$$ as $$|0\rangle = |0000\rangle$$. This error is illustrated by the green square in the right panel of Fig. 2(b), which also shows the measurement outcomes when a single ancilla readout or storage decay error occurs.

### IV. MEASUREMENT ERROR MITIGATION

Can we use our knowledge of the error mechanisms described above to improve the fidelity of our measurements? To answer this question, we use an approach known as error mitigation, which relies on data postprocessing to improve the quality of an ensemble of calculations or measurements. Error mitigation has been proposed [53,54] and realized [55] for use with variational algorithms. Here, our error mitigation efforts...
focus on an ensemble of measurements, which sample from the population distribution of some quantum state.

In this section, we introduce an error model based on the error syndromes described in Sec. III C. This model is a function of several error rates, which must be calibrated. We discuss our calibration technique and compare the model to experimental results. Finally, we introduce the deconvolution techniques that use the calibrated model to mitigate the errors in the measurement and show an improvement on two sets of states.

A. Error model

Our goal is to characterize the errors so that we can model measurement outcomes. This is equivalent to finding the elements \( F_{i} \) of the positive-operator-valued measure (POVM) that describes our measurement [56,57]. For instance, the bitwise measurement is composed of \( B \) individual bit measurements, each with a corresponding POVM \( E_{b_{j} \in \{0,1\}}^{(k)} \) where \( E_{b_{j}}^{(k)} = (\hat{B}_{b_{j}}^{(k)})^\dagger \hat{B}_{b_{j}}^{(k)} \). Each ideal \( F_{i} \) is then a product of POVM elements constituting the bitwise measurement:

\[
F_{i=b_{a-1}...b_{0}} = E_{b_{a-1}}^{(B-1)} \cdots E_{b_{0}}^{(0)}.
\] (7)

Note that the \( F_{i} \) and \( E_{b_{j}}^{(k)} \) are diagonal because the bit measurement operators \( \hat{B}_{b_{j}}^{(k)} \) are QND. Ideally, \( F_{i} = |i \rangle \langle i | \), but our implementations of the \( E_{b_{j}}^{(k)} \) have errors, so \( (F_{i})_{j} \) is the probability that the measurement detects \( i \) photons when there are \( j \) photons in the mode \( P(\text{outcome } = i | j) \).

Our system can be described with a hidden Markov model (HMM) that has been used for qubit readout [58–62] and to improve the readout of qubits encoded in oscillators [51,63]. A HMM parametrizes the errors in the storage mode, known as transitions, and the fidelity of the measurements of each bit, known as emissions. The probability of a transition in the storage mode during a time interval \( t \) from \( |i \rangle \) to \( |j \rangle \) is

\[
T_{i,j}(\kappa t) = \left[ \left( \begin{array}{c} i \\ j \end{array} \right) \right] (e^{\kappa t} - 1)^{j-i} e^{-i\kappa t}
\] (8)

where \( \kappa \) is the decay rate of the mode [63]. The ancilla can decay during the controlled rotation \( C\hat{P}^{k} \) and the readout procedure, leading to an incorrect bit measurement. Since the bit measurements are QND, only the diagonal entries \( E_{b_{j}}^{(k)} \) are nonzero, even in the presence of errors. We modify the entries of these ideal POVM elements

\[
E_{b_{j},i}^{(k)} = \left\{ \begin{array}{ll}
\frac{i}{j} & \text{is even} \\
\frac{j}{i} & \text{is odd}
\end{array} \right.
\]

(9)

\[
to introduce error rates \( \epsilon_{(k)j} \) describing errors in the \( C\hat{P}^{k} \) mapping or ancilla readout.

We now integrate both error mechanisms into a single model describing the errors in the measurement. The bitwise measurement algorithm described in Sec. III C alternates between entangling pulses and ancilla reset. The cavity is subject to spontaneous decay for the duration of the algorithm, punctuated by projective measurements of the entangled ancilla. This duration, however, is not deterministic due to the dynamic ancilla reset. Our HMM models the errors in the measurement by constructing an alternating chain of transition and emission (measurement) events. To calculate the \( F_{i} \), we sum over all transition paths \( s_{k} \) between storage state \( |j \rangle \) and measurement outcome \( i \). Weighting these paths by the probability that the bit measurements yield the binary decomposition of \( i = b_{1}b_{2}b_{3}b_{0} \) gives

\[
(F_{i})_{j} = \sum_{s_{1}...s_{B}} T_{s_{1},s_{1}}(\kappa t^{(0)}) E_{b_{0},s_{1}}^{(0)} T_{s_{1},s_{1}}(\kappa t^{(1)} + b_{0}\kappa t') \times E_{b_{1},s_{2}} \cdots E_{s_{B-1},s_{B}}(\kappa t^{(B-2)} + b_{B-1}\kappa t') E_{b_{B-1},s_{B}}^{(B-1)}.
\] (10)
Each transition matrix $T$ and POVM element $E^{(k)}_{i,j}$ are parametrized by the respective error rates $\kappa_t^{(k)}$ and $\epsilon_{g}^{(k)}$, $\epsilon_{s}^{(k)}$, which are not necessarily the same for each bit. Note that the time between each bit measurement depends on each result, as a measurement of $b_k = 1$ requires an ancilla reset pulse followed by readout verification. We discuss how the dynamic ancilla reset statistics inform our choice of $\kappa_t$ in Appendix A. Equation (10) provides the desired parametrized POVM elements, which, once calibrated, we use to mitigate errors in our measurement.

B. Model calibration and results

Now that we have a model for the errors, we must determine the values of the error parameters $\kappa_t^{(k)}$ and $\epsilon_{g}^{(k)}$, $\epsilon_{s}^{(k)}$. The duration of each bit measurement depends on each result, as a measurement of $b_k = 1$ requires an ancilla reset pulse followed by readout verification. We discuss how the dynamic ancilla reset statistics inform our choice of $\kappa_t$ in Appendix A. Equation (10) provides the desired parametrized POVM elements, which, once calibrated, we use to mitigate errors in our measurement.

Once we have calibrated the error rates in the model (tabulated in Appendix C), we use Eq. (10) to calculate the $\{F_i\}$ and thus reconstruct the detector POVM. We compare the modeled POVM to the measured POVM in Fig. 3 by preparing the Fock states $|i\rangle_{b_6,j_6,15}$ and measuring all four bits. The residuals are typically 1%, and never more than 4%, showing good agreement with the model and bounding unmodeled errors. We could in principle use the measured detector POVM to perform error mitigation, but we would still need to model transitions during the Fock state preparation and would not be able to confirm our understanding of the error mechanisms.

The raw measurement fidelity of each state is at least 80% before error mitigation even for the most challenging input state $n = 15$. The measurement extracts on average between 3.11 and 3.17 of the four possible bits of information (the range exists due to state preparation details; see Appendix E for further discussion). This fraction is greater 3/4, suggesting that the bitwise measurement is single shot.

C. Error deconvolution

Now that we have modeled the entire collection of $\{F_i\}$, we form a matrix $C_{ij} = (F_i)_{j_6}$, where $C = I$ is shown in Fig. 3(b), of conditional probabilities referred to as the “confusion matrix.” This allows us to relate the measured populations $P_{\text{meas}}$ to the state’s ideal distribution $P_{\text{ideal}}$ with a simple relation [57]

$$P_{\text{meas}} = CP_{\text{ideal}}. \tag{11}$$

Once we have $C$, we invert it to solve the above equation for the mitigated probability vectors $P_{\text{mit}} = C^{-1}P_{\text{meas}}$ (Appendix F contains a discussion of the invertibility of $C$). These vectors are properly normalized, but may have negative entries. We correct the mitigation results by finding the closest vector $\tilde{P}_{\text{mit}}$ with respect to the Euclidean norm that has non-negative entries [57].
single bit error rates \( \epsilon \) compared to the error in the mitigated measurements \( D_{TV}(P_{\text{mit}}, P_{\text{ideal}}) \) (blue) is compared to the error in the mitigated measurements \( D_{TV}(P_{\text{mit}}, P_{\text{ideal}}) \) (orange). The large difference in overall error rates between Fock and coherent states results from the error mechanism to which each is most susceptible. Fock states transition to orthogonal states under photon loss, while coherent states are mainly affected by bit measurement errors. The single photon loss probability \( P_0 \) is a high-fidelity set of states, and directly inverting the error model, we apply it to the same Fock states we prepared in the calibration process, and to a set of coherent states with 0 \( \leq \alpha \leq 3 \). The results are shown in Fig. 4. However, it turns out that calculating any particular element of \( P_{\text{mit}} \) from \( P_{\text{meas}} \) is efficient, as we show below. Calculating particular elements of \( P_{\text{mit}} \) may be sufficient for problems in which the goal is to study features of an output distribution that can be altered by measurement errors. For instance, we may be interested in calculating the relative peak intensities of Franck-Condon factors \([24,49]\), which only transitions to an orthogonal state, leaving bit measurement errors as the dominant error mechanism.

The error mitigation successfully reduces errors in all tested states to less than 5\%, providing the most dramatic improvement to the Fock states, whose average error is reduced from 13.5\%–1.1\%. The success of the error mitigation shows that we understand the error syndromes in our measurement and can effectively quantify them. This allows us to use the model to find the error rates required to meet a desired measurement fidelity, which may vary from application to application.

D. Scalability of error mitigation

In the previous sections, we demonstrated an efficient measurement of the number of photons in a mode, and an effective method to mitigate the errors in the measurement. Our protocol can be performed simultaneously on a system with multiple storage modes, each dispersively coupled to their own ancilla qubit, to measure the joint photon number of the system. Imagine a system with \( M \) storage modes with uncorrelated measurement errors between the modes. Here, the entries of \( P_{\text{meas}} \) are elements of the joint probability distribution \( P(n_1, n_2, \ldots, n_M) \). If we truncate the Hilbert space of each mode to some maximal number of photons \( N_{\text{max}} \), then the length of \( P_{\text{meas}} \) is \( N_{\text{max}}^M \). This exponentially large space suggests that our approach to error mitigation is not scalable to multiple modes.

The simple form of Eq. (11) offers the tantalizing possibility that we could avoid error modeling altogether by preparing a high-fidelity set of states, and directly inverting \( P_{\text{ideal}} \) to find \( C \). Coherent states seem to be an ideal candidate due to the high fidelity with which we can prepare them, but suffer from ill-conditioned ideal populations \( P_{\text{ideal}} \), which make the determination of \( C \) highly sensitive to errors in \( P_{\text{meas}} \) [65]. Furthermore, error modeling of the Fock states is required due to preparation error discussed in the previous section.

To quantify the quality of our measurement, we use the total variation distance metric to compare probability distributions \( A \) and \( B \),

\[
D_{TV}(A, B) = \frac{1}{2} \sum_k |A_k - B_k|. \tag{12}
\]

The factor of 1/2 accounts for the double counting of probability differences.

To check the performance of the error mitigation and verify the error model, we apply it to the same Fock states we prepared in the calibration process, and to a set of coherent states with 0 \( \leq \alpha \leq 3 \). The results are shown in Fig. 4. In both cases, the error mitigation includes photon losses occurring during state preparation (recall the selective pulse check discussed in Sec. IV B). This is reasonable because the operation that directly precedes the bitwise measurement is known in advance as part of the pulse sequence.

The large difference in the premitigated errors in Fig. 4 between the Fock and coherent states is due to the nature of these states. For the Fock states, a single lost photon transitions [\( n \)] to an orthogonal state [\( m \)], resulting in the maximum possible error. Furthermore, the transition rate increases with photon number, resulting in large premitigation errors. On the other hand, coherent states are more robust to photon loss and do not transition to an orthogonal state, leaving bit measurement errors as the dominant error mechanism.

FIG. 4. Error mitigation results for coherent (left) and Fock states (right). The \( \alpha \) were chosen such that the Hilbert space can be truncated to the first 16 Fock states, with error less than 2.2\%. For each prepared state [\( |\psi\rangle \)], the error in the measurements \( D_{TV}(P_{\text{mit}}, P_{\text{ideal}}) \) (blue) is compared to the error in the mitigated measurements \( D_{TV}(P_{\text{mit}}, P_{\text{ideal}}) \) (orange). The large difference in overall error rates between Fock and coherent states results from the error mechanism to which each is most susceptible. Fock states transition to orthogonal states under photon loss, while coherent states are mainly affected by bit measurement errors. The single photon loss probability \( \kappa t^{(k)} \approx 0.03 \) is similar to the single bit error rates \( \epsilon_k^{(k)} \approx 0.03 \), \( \epsilon_k^{(k)} \approx 0.01 \).
requires processing the measured peaks, not the entire Hilbert space. The resulting spectrum will have significant peaks at the same output photon numbers, with intensities adjusted to account for measurement errors, increasing the simulator’s accuracy. This way, the simulator still does all the hard work of identifying the significant peaks; we only need to postprocess the heights to mitigate measurement errors.

There are two properties that make calculating any element of $P_{\text{mit}}$ efficient. The first is that the multimode confusion matrix is the Kronecker product of single-mode confusion matrices, meaning that it is efficient to calculate any element in the inverse of the multimode confusion matrix. In general, though, $C^{-1}$ contains an exponential number of nonzero matrix elements so writing down the entire matrix is not feasible. The second property is that the vector $P_{\text{mit}}$ is necessarily sparse; assuming we do not perform an exponentially large number of measurements, most configurations $(n_1, n_2, \ldots, n_M)$ are never actually measured. Together, these two properties imply that any entry of $P_{\text{mit}}$ can be efficiently computed,

$$P_{\text{mit}} = \sum_{j \in S} (C_1^{-1} \otimes C_2^{-1} \cdots \otimes C_M^{-1})_{ij} P_{\text{meas}}^{\text{mit}},$$

where $S$ is the set of nonzero entries in $P_{\text{meas}}$, which is only polynomially large. We cannot use this fact to efficiently calculate the entire distribution $P_{\text{mit}}$, however, because $P_{\text{mit}}$ can contain exponentially many nonzero entries. In the case where simply adjusting peak intensities is insufficient, we have developed a method for expanding $C^{-1}$ that mitigates errors up to a chosen order. This method is detailed in Appendix G.

V. CONCLUSION AND OUTLOOK

Our bitwise measurement protocol utilizes the long lifetimes of three-dimensional (3D) microwave cavities and the interactions generated by a Josephson junction to implement a high-fidelity, single-shot, microwave photon-number-resolving detector. We use error mitigation techniques to improve the measurement fidelity of the Fock states by nearly an order of magnitude. Larger resolvable photon numbers are reachable by synthesizing additional $C^{-1}$ pulses to measure more bits.

In addition to enabling error mitigation, our error model provides insight into the error budget of our PNR detector. For example, we estimate that with unit fidelity $C^{-1}$ pulses and ancilla readout ($\epsilon^{(k)} = \epsilon^{(k)} = 0$) we can extract 3.72 bits of information in a single shot in our system with the current $\kappa T(k)$. This suggests that further optimization of the control pulses and ancilla readout can increase the information extracted per shot by up to 0.6 bits. Any additional improvements of the single-shot fidelity require decreasing $\kappa T(k)$ by decreasing the pulse lengths or increasing cavity lifetimes. Adaptive check methods may also be able to improve single-shot fidelity at the expense of adding complexity to the error model.

Even with limited fidelity, optical PNR detectors have motivated the development of multiphoton quantum information processing protocols. Bringing high-fidelity PNR detectors to the microwave regime further bolsters the ability of cQED to perform photonic protocols, such as vibronic spectra simulations [49]. Additionally, this capability will motivate the development of algorithms that take advantage of the large bosonic Hilbert space available in cQED systems.

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APPENDIX A: DYNAMIC ANCILLA RESET STATISTICS

As discussed in the Supplemental Material in Ref. [51], our ancilla qubit reset protocol does not always succeed in a single attempt. This increases the duration between bits on a shot-by-shot basis. This duration is proportional to the number of reset attempts, whose probabilities shown in Fig. 5 have a long tail, with 35.1% of resets requiring more than one attempt. The average number of reset attempts required across all Fock states prepared in Figs. 3 and 4(b) is $\bar{n}_{\text{reset}} = 2.05$. Each reset attempt has a duration of 2.244 $\mu s$ so we use the average reset time $\kappa T' = 2.244 \mu s \cdot \kappa \bar{n}_{\text{reset}} = 0.0046$ in the model.

It is tempting to remove runs via postselection that take more than a chosen threshold of reset attempts, but this introduces bias into the measurement results as the number of required reset attempts increases with the number of photons in the storage mode. Postselecting on a low threshold such as five attempts (used in Ref. [51]) removes 5% more...
TABLE I. HMM error parameters. The error $\sigma$ of $\epsilon^{(k)}$ and $\epsilon^{(j)}$ is calculated by propagating error in Eq. (B2). 1/$\kappa$ $\approx$ 1 ms is the storage cavity decay rate.

<table>
<thead>
<tr>
<th>Bit $k$</th>
<th>$\kappa t^{(k)}$ $(+/\kappa t^{(j)})$</th>
<th>$\epsilon^{(k)}$ $\pm \sigma$</th>
<th>$\epsilon^{(j)}$ $\pm \sigma$</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>0.0040, 0.0032$^a$</td>
<td>0.019 $\pm$ 0.002</td>
<td>0.029 $\pm$ 0.002</td>
</tr>
<tr>
<td>1</td>
<td>0.0034 $(+/0.0046)^b$</td>
<td>0.014 $\pm$ 0.001</td>
<td>0.026 $\pm$ 0.001</td>
</tr>
<tr>
<td>2</td>
<td>0.0034 $(+/0.0046)^b$</td>
<td>0.011 $\pm$ 0.001</td>
<td>0.035 $\pm$ 0.002</td>
</tr>
<tr>
<td>3</td>
<td>0.0034 $(+/0.0046)^b$</td>
<td>0.013 $\pm$ 0.001</td>
<td>0.033 $\pm$ 0.002</td>
</tr>
</tbody>
</table>

$^a$Fock and coherent states, respectively, which have different preparation methods.

$^b$The parenthetical additional time is added when the ancilla needs to be reset and is equal to $\kappa t^{(j)}$, as discussed in Appendix A.

measurements of 15 than 1. Setting the postselection threshold at 50 attempts lowers this disparity to 0.6%, but we still choose to avoid postselection to produce an unbiased measurement. A more granular approach to ancilla reset using the state-dependent reset probability distributions can be included in the HMM if one separately calibrates this dependence.

APPENDIX B: CHOICE OF CALIBRATION STATES

Consider a set of calibration states $\{|\psi_j\rangle\}$. After preparing each $|\psi_j\rangle$, we measure the $k$th bit to find the probability of measuring $b_j = 0$, 1 for each basis state $P_{b_{j,k}} = P(b_j = 0, 1 | \psi_j\rangle) = \langle \psi_j | E^{(k)}_{b_j} | \psi_j\rangle$. Our preparation protocol may suffer from transitions caused by ancilla readout and storage mode decay that we can model in the calibration process

$$P_{b_{j,k}} = \sum_{m,n} T_{m,n}(\kappa t^{(k)}) E^{(k)}_{b_j,n} O_{j,m}, \quad (B1)$$

where $O_{j,m} = \langle \psi_j | m \rangle^2$ changes from the calibration basis $\{|\psi_j\rangle\}$ to the Fock basis. The argument of $T$, $\kappa t^{(0)}$, includes the duration of the final check of the preparation protocol and the $C^{b_{j,k}}$ pulse. It can be shown that the effective transition duration for this final check is approximately half of the length of the selective pulse. The value of $\kappa t^{(0)}$ is given in Table I. Solving for the diagonal of $E^{(k)}_{b_j}$ we find

$$E^{(k)}_{b_j} = (T^{-1}(\kappa t^{(0)})O^{-1}P_{\text{cal}})_{j,j}, \quad (B2)$$

from which we average and extract $\epsilon^{(k)}$, $\epsilon^{(j)}$.

Our task is then to pick a basis of calibration states $|\psi_j\rangle$. In our systems, we can prepare coherent states $|\alpha\rangle$ with the highest fidelity due to the speed of preparation. If we choose a calibration basis $|\alpha_j\rangle$ such that there is very little probability finding $n$ photons with $n > N_{\text{max}}$, we would expect these states to be an excellent candidate for a calibration basis. However, the presence of $O^{-1}$ in Eq. (B2) complicates matters. To satisfy $n \leq N_{\text{max}}$, the range of $\alpha_j$ is necessarily restricted. In this regime, the coherent states are only barely linearly independent and have significant overlap $|\alpha_j|^{2}$. This results in $O$ being an ill-conditioned matrix, meaning that $O^{-1}$ magnifies small errors in $P_{\text{cal}}$. We avoid the ill-conditioning problem by using the Fock states $|j\rangle$ as our calibration basis resulting in $O = 1$.

APPENDIX C: MODEL ERROR RATES

Table I contains error parameters calibrated in Sec. IV B and used in the HMM.

APPENDIX D: $C^{b_{j,k}}$ PULSE CONSTRUCTION

Optimal control techniques are powerful tools to engineer quantum gates and initialize nonclassical quantum states. We use this tool here to create the $C^{b_{j,k}}$ entangling gate in a manner very similar to that described for a cavity-transmon system in Ref. [41].

In particular, we use gradient ascent pulse engineering (GRAPE) to compute both quadratures of an ancilla control pulse that realizes $C^{b_{j,k}}$. The optimizer uses the cavity-qubit Hamiltonian in the rotating frame

$$\hat{H}_{\text{GRAPE}}/\hbar = -\frac{\chi_{\sigma}}{2} a^\dagger a^\dagger a - \frac{K}{2} s^\dagger s^\dagger s \hat{s}^\dagger \hat{s}^\dagger \hat{s}^\dagger \hat{s}^\dagger a - \frac{\chi_{\sigma}}{2} s^\dagger \hat{s}^\dagger \hat{s}^\dagger \hat{s}^\dagger a + I(t)^{2}(\hat{a} + \hat{\alpha}) + iQ(t)(\hat{a} - \hat{\alpha}) \quad (D1)$$

to compute the best $\Omega(t) = I(t) + iQ(t)$ that implements $C^{b_{j,k}}$. We convert $\Omega(t)$ into a DAC amplitude with a scaling factor determined by calibrating the Rabi rate of a $\pi$ pulse on the ancilla qubit. The substitution $\hat{\alpha} \rightarrow \hat{a} \hat{\alpha}$ allows GRAPE to optimize over multiple transmon levels. We do not include loss in the qubit and cavity in the optimization. The optimizer finds the ideal control pulse $\Omega$ by maximizing the fidelity of the state transfer in Eq. (4)

$$\mathcal{F}(\Omega) = \sum_k \left|\langle \psi^{(k)}_{\text{initial}} | \hat{U}(\Omega) | \psi^{(k)}_{\text{target}} \rangle\right|^2, \quad (D2)$$

where $\hat{U}(\Omega)$ is the unitary implemented by the pulse $\Omega$. Note that the modulus is taken within the sum, so the optimal pulse produced by the optimization maximizes the incoherent fidelity of the $C^{b_{j,k}}$ gates. Both quadratures of these pulses are shown in Fig. 6 with Hamiltonian parameters listed in Table II. Each pulse reaches a fidelity of at least 99.9% in a lossless system.

APPENDIX E: CALCULATION OF INFORMATION EXTRACTED BY THE DETECTOR

To be able to call our detector single shot, it must extract a large fraction of the information it is designed to measure per shot. In particular, of the four bits it ideally measures, we would like to determine how many of these bits are lost to measurement errors. The difference of ideal and wasted bits gives the average number of bits extracted per measurement. Our task is then to calculate the number of wasted bits in our measurement, which is equivalent to number of extra bits needed to classify the input given a measurement result (or the number of bits remaining to be measured). This quantity is the entropy $S_i$ of the distribution $P(|j\rangle | i) = P(\text{detector input } | j\rangle | \text{measurement result } i)$. We finally average over the results $i$ to obtain $4 - (\bar{S}_i)$, the average number of bits extracted in a single shot of the measurement.
We begin by calculating the conditional probability

\[ P(j|i) = \frac{P(\text{result } j|\text{ input } j)P(|j\rangle)}{P(i)} \]  

(E1)

Recalling that \( P(\text{result } j|\text{ input } j) = C_{ij} \) and assuming a uniform prior \( P(j|i) = 1/N \) we simplify Eq. (E1)

\[ P(j|i) = \frac{C_{ij}P(|j\rangle)}{\sum_k C_{ik}P(|k\rangle)} = \frac{C_{ij}}{\sum_k C_{ik}}. \]  

(E2)

Finally, we write down the average entropy over all measurement outcomes, weighting by \( P(\text{outcome } i) \)

\[ \langle S_i \rangle = -\sum_{ij} P(|j\rangle|i) \log_2(P(|j\rangle|i))P(i) \]

\[ = -\frac{1}{N} \sum_{ij} C_{ij} \log_2 \left( \frac{C_{ij}}{\sum_k C_{ik}} \right). \]  

(E3)

Using the error parameters listed in Table I to construct \( C \), we find \( 4 - \langle S_i \rangle_{\text{Fock}} = 3.14 \) bits. We can also use the measurement results in Fig. 3 to determine \( P(\text{result } j|\text{ input } j) \), yielding \( 4 - \langle S_i \rangle_{\text{Fock}} = 3.11 \) bits. This number likely provides an upper bound of \( \langle S_i \rangle \) as the check protocol for Fock states (described in Sec. IV B) introduces a lengthy delay between state preparation and measurement.

As mentioned above, the preparation errors included in both of these calculations likely inflate the number of bits the measurement itself wastes. We cannot deconvolve only the preparation error from the measurement results, but we can zero the contribution of preparation to \( \kappa^{(0)} \). We emphasize that \( \kappa^{(0)} \) will always contain a contribution from the previous experimental step, but we can construct a \( C \) confined to errors that solely occur during the measurement duration. This scenario yields \( 4 - \langle S_i \rangle_{\text{no prep}} = 3.17 \) bits, which is likely a lower bound as the residuals in Fig. 3 show that \( C \) does not perfectly capture all errors.

Finally, we can use this method to estimate the single-shot error budget. With \( \kappa^{(k)} \) as in Table I and unit fidelity \( \hat{C}^{(k)} \) pulses and ancilla readout (\( \epsilon_g^{(k)} = \epsilon_e^{(k)} = 0 \)) we find \( 4 - \langle S_i \rangle_{\text{only } \kappa^{(k)}} = 3.72 \) bits. The remaining information is lost to photon loss in the storage mode. This suggests that significant gains can be made by optimizing the \( \hat{C}^{(k)} \) pulses and ancilla readout before needing to decrease \( \kappa^{(k)} \).

### APPENDIX F: INVERTIBILITY OF \( C \)

Our error mitigation protocol works in the range of error rates that allow robust inversion of \( C \). In the limit of small error rates, \( C \) approaches the identity. As the error rates \( \epsilon_g^{(k)}, \epsilon_e^{(k)} \), and \( \kappa^{(k)} \) become large, we expect that \( C \) will become ill-conditioned. For example, as the storage mode error rates \( \kappa^{(k)} \) approach 1, the probability \( C_{0n} \) of any photon state \( |n\rangle \) decaying to \( |0\rangle \) approaches 1. But since the column vectors of \( C \) are normalized to 1, they approach degeneracy and prevent the inversion of \( C \). This exact scenario is illustrated in Fig. 7. Similarly, the condition number also diverges as \( \epsilon_g^{(k)}, \epsilon_e^{(k)} \to 1 \). We operate the small error rate regime in which \( C \) is robustly invertible.

### APPENDIX G: CONFUSION MATRIX EXPANSION

If simply correcting peak intensities in a measured spectrum is insufficient, we can expand the confusion matrix about the identity matrix. By truncating this expansion to some
order so that the total number of relevant matrix elements is polynomial, we can calculate polynomially many of the largest entries in $P_{\text{ideal}}$. Using this approach, we may find significant peaks $P_{\text{ideal}} \neq 0$ even when $P_{\text{meas}} = 0$. However, each peak height will only be accurate to the order at which the expansion is truncated.

Consider the confusion matrix $C_m$ for photon number read-out of the $m$th mode. Let us assume that $C_m$ is close to the identity, so that its inverse is also close to the identity,

$$C_m^{-1} = I + \varepsilon E_m,$$

where $\varepsilon \ll 1$ is some small number, and $E_m$ is a matrix normalized so that

$$\|E_m\|_{\text{max}} = \max_{ij} |E_m|_{ij} = 1.$$  \hspace{1cm} (G1)

Thus,

$$P_{\text{meas}} = [(I + \varepsilon E_1) \otimes (I + \varepsilon E_2) \otimes \ldots \otimes (I + \varepsilon E_M)]P_{\text{meas}}.$$  \hspace{1cm} (G3)

Now, the main idea is to use this expansion to compute only the entries of $P_{\text{meas}}$, which are $O(\varepsilon^q)$, then neglect all entries, which go as higher powers in $\varepsilon$. To perform this computation, first notice that we only need to consider a small subset of the columns of $C_m^{-1}$ because $P_{\text{meas}}$ is sparse. That is, only the columns $[C_m^{-1}]_{\varepsilon, j \in S}$ contribute, where $[C_m^{-1}]_{\varepsilon, j}$ denotes the $j$th column of $C_m^{-1}$. While each such column contains exponentially many nonzero entries $N^{\text{max}}_m$ in general, it turns out that the columns may contain only polynomially many nonzero entries when expanded in $\varepsilon$, depending on the order of the expansion. To compute $P_{\text{meas}}$ we then compute all entries, which are nonzero to order $q$ in the columns $[C_m^{-1}]_{\varepsilon, j \in S}$, giving us a sparse representation of $C^{-1}$. It is then efficient to multiply $C^{-1}P_{\text{meas}}$.

Each entry in $C^{-1}$ is a product of entries from the single-mode confusion matrices $C^{-1}_k$. In each column of $C^{-1}_k$, there are $N^{\text{max}}_m$ elements, one of which is $O(1)$ and up to $N^{\text{max}}_m - 1$ of which are $O(\varepsilon)$ as defined in Eq. (G1). Each of the $N^{\text{max}}_m$ elements in a column of $C^{-1}$ is a product of $M$ elements where the $k$th element in the product is drawn from a fixed column of $C^{-1}_k$. Our task is to count the number such products that result in an element of $O(\varepsilon^q)$. We choose $q$ columns out of $M$ to each contribute a factor of $\varepsilon$, and within each column we choose one of $N^{\text{max}}_m - 1$ entries that are $O(\varepsilon)$. There are thus at most $\binom{M}{q}(N^{\text{max}}_m - 1)^q$ entries in a column of $C^{-1}$, which are polynomials of degree $q$ in $\varepsilon$. Therefore, the total number of nonzero entries $N$ per column, which must be calculated when expanding $[C^{-1}]_{\varepsilon, j \in S}$ to $q$th order in $\varepsilon$ is

$$N = \sum_{\ell=0}^q (N^{\text{max}}_m - 1)^q \binom{M}{\ell}.$$  \hspace{1cm} (G2)

If $N$ is only polynomially large, this approach is efficient, as the sparsity of $P_{\text{meas}}$ ensures that the total number of entries $N[S]$ in $[C^{-1}]_{\varepsilon, j \in S}$ is not exponentially large.

We now examine the scaling of $N$, which can be written as $N = O(N^{\text{max}})^q \times \sum_{\ell=0}^q \binom{M}{\ell}$. The partial sum of binomial coefficients $\sum_{\ell=0}^q \binom{M}{\ell}$ does not have a closed form expression; for small values of $q$ it is $O(M^q)$, but for $q = M$ it is $O(2^M)$. Thus, whether this approach is efficient depends on the choice of $q$. For small $q$, we have $N = O(N^{\text{max}}_m M^q)$ so that the total number of nonzero entries scales polynomially in the number of modes $M$ and the maximal photon number $N^{\text{max}}_m$, ensuring the efficiency of the expansion.


[57] F. B. Maciejewski, Z. Zimborás, and M. Ozsmianeck, Mitigation of readout noise in near-term quantum devices by classical post-processing based on detector tomography, Quantum 4, 257 (2020).


